

Sufficient Conditions for Robustness of RDMA Programs

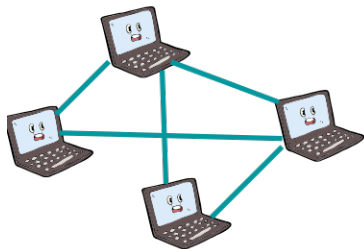
ESOP 2025

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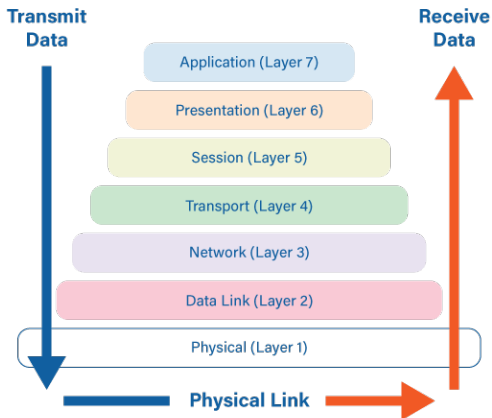
May 8, 2025

HPC/database: slow communication

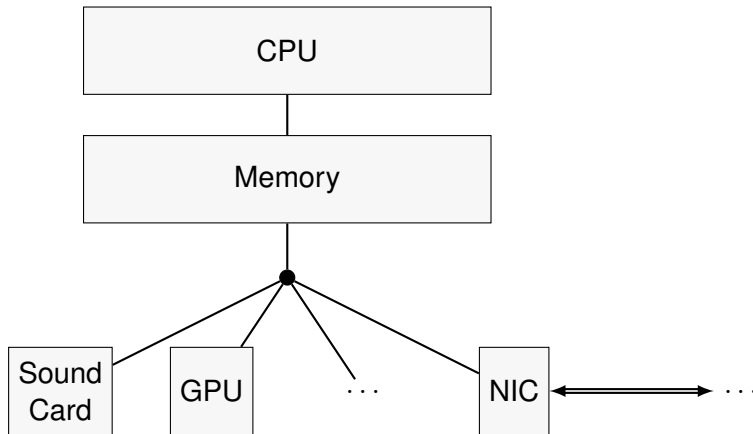


Latency : $\sim$ ms

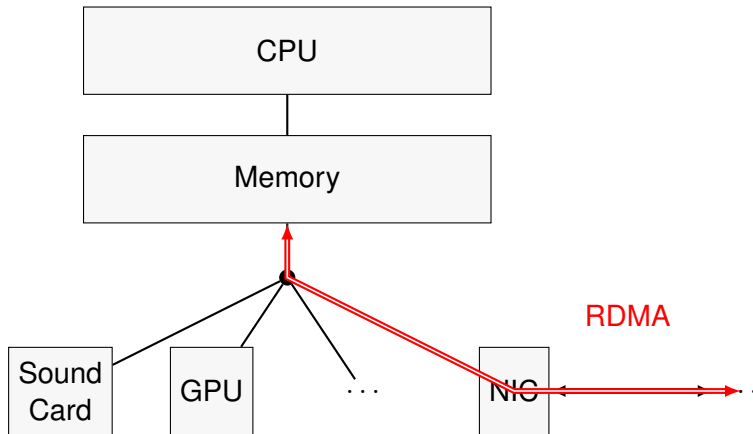
The 7 Layers of OSI



Direct Memory Access



Direct Memory Access



Remote Direct Memory Access

RDMA

- Data-transfer protocol for High-Performance Computing
- Latency : $\sim \mu s$
- Recently becoming widespread
- Very weak and unintuitive semantics¹

Our work: robustness as a path for verification

- Making an existing program robust
- Guidelines for writing robust programs

¹Ambal et al., OOPSLA 2024

Remote Direct Memory Access

RDMA

- Data-transfer protocol for High-Performance Computing
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¹Ambal et al., OOPSLA 2024

Model and Syntax

Model:

- network with several nodes ($n_1, n_2, \dots$)
- each node can have several threads

Syntax

Mem Locs $x, y, z, \dots$

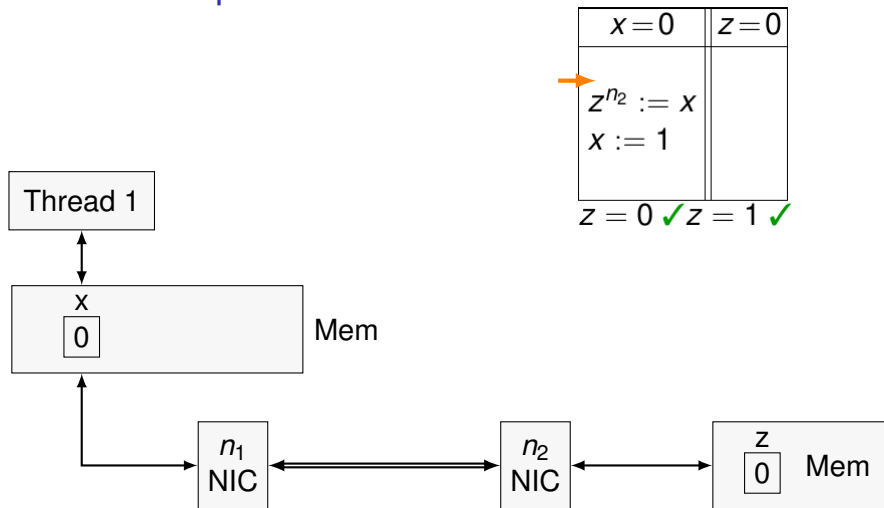
Expressions $e ::= k \in \mathbb{Z} \mid x \mid e_1 + e_2 \mid \dots$

Commands $c ::= x := e \mid \text{mfence} \quad \leftarrow \text{usual x86-TSO}$
 $\quad \mid x^n := y \mid x := y^n \mid \text{poll}(n) \quad \leftarrow \text{RDMA}$

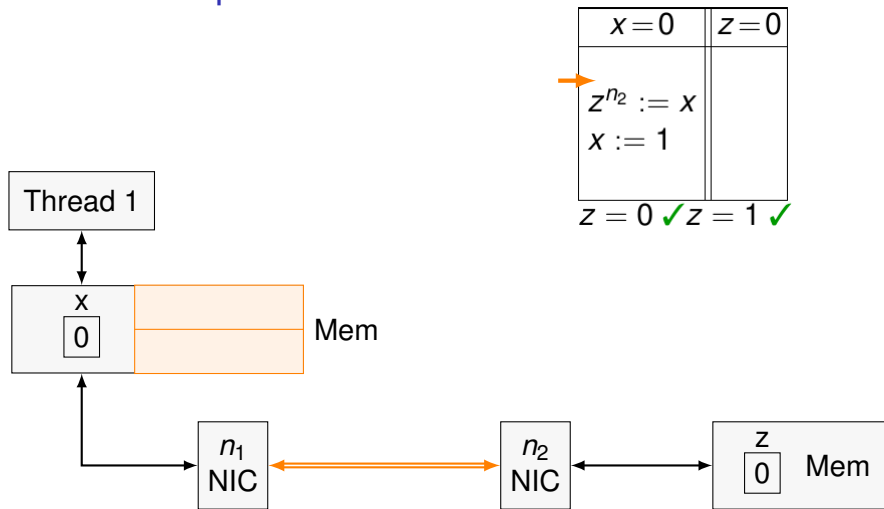
RDMA example 1/3

$x = 0$	$z = 0$
$z^{n_2} := x$ $x := 1$	
$z = 0$ ✓	$z = 1$ ✓

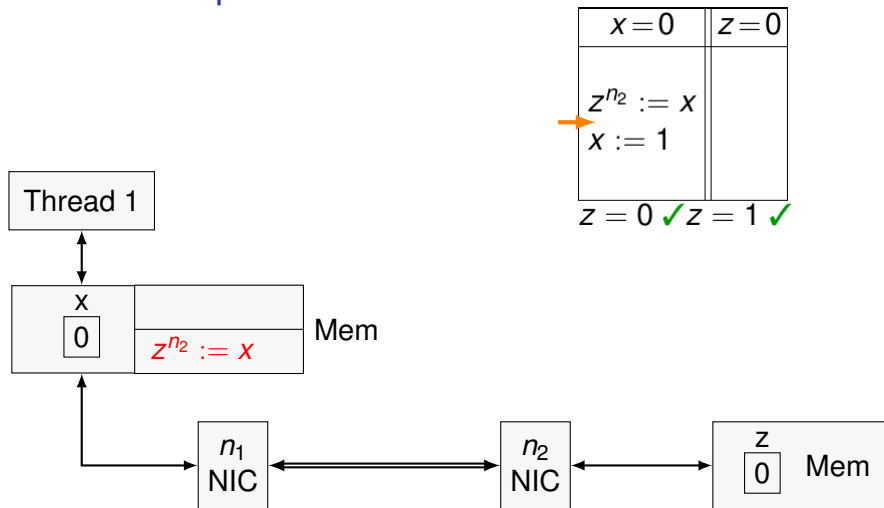
RDMA example 1/3



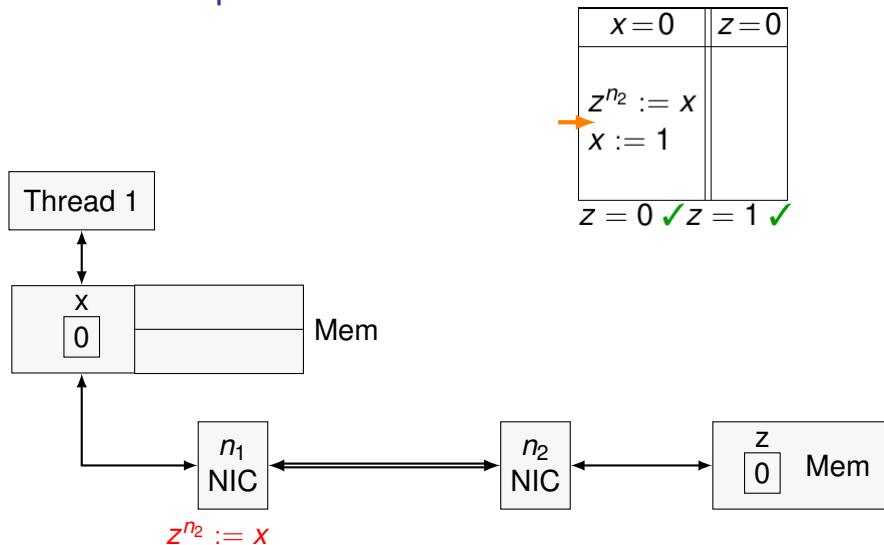
RDMA example 1/3



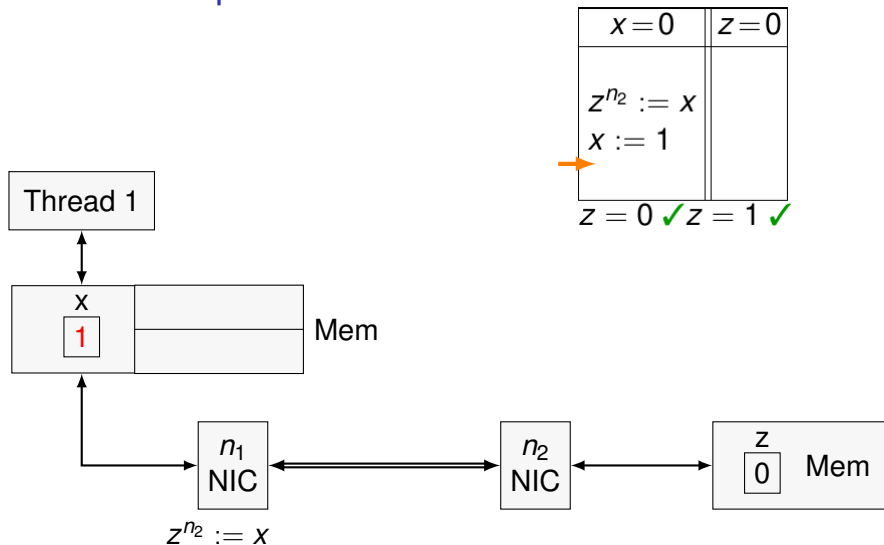
RDMA example 1/3



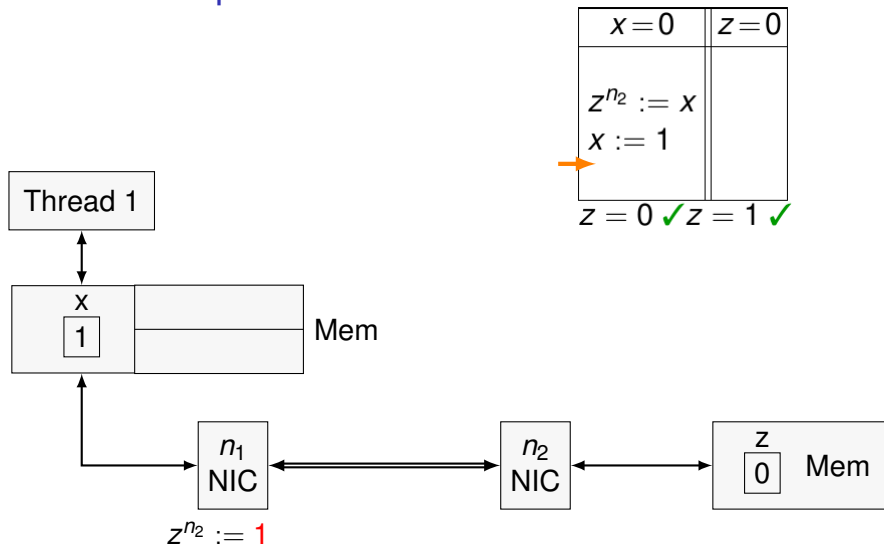
RDMA example 1/3



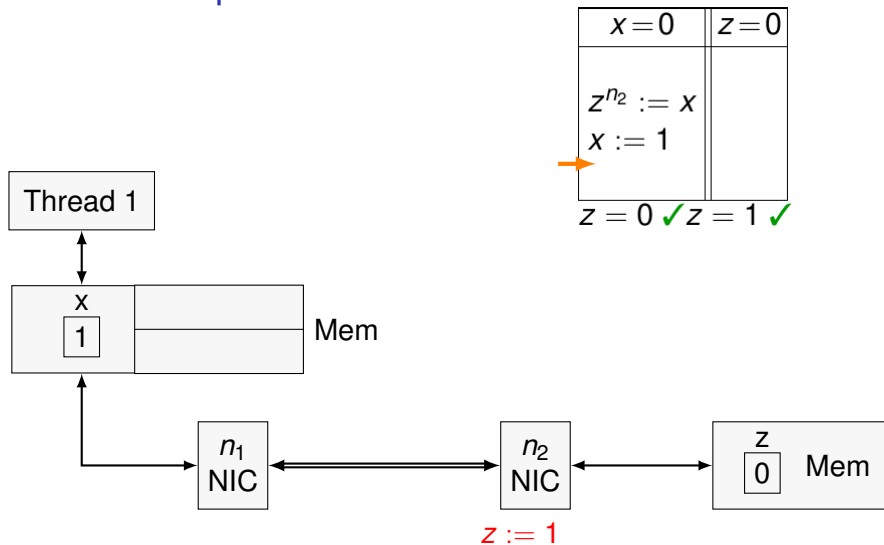
RDMA example 1/3



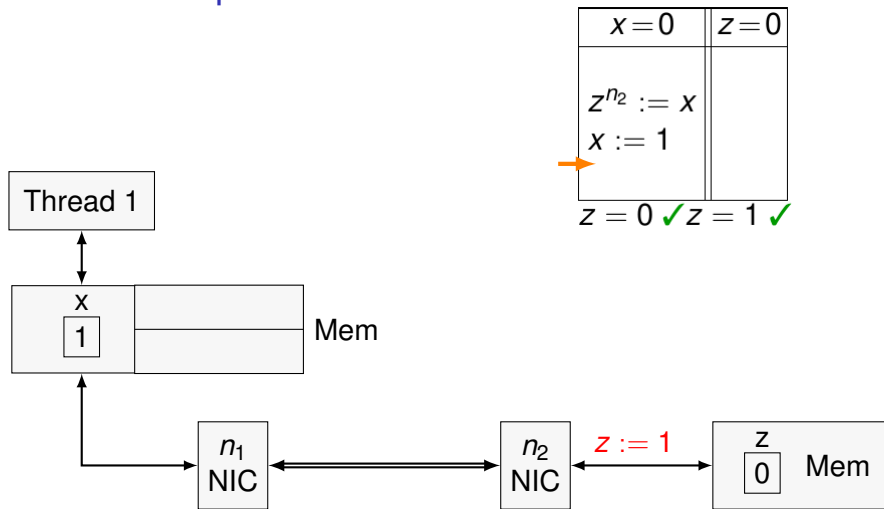
RDMA example 1/3



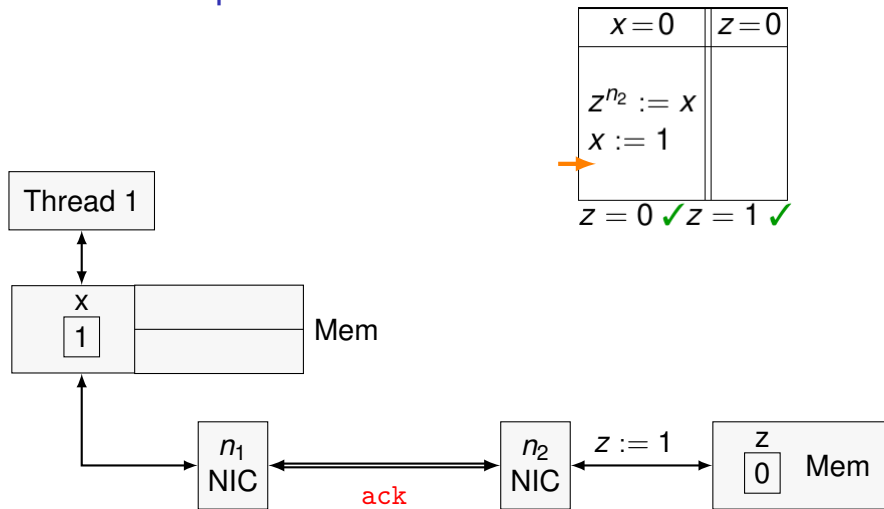
RDMA example 1/3



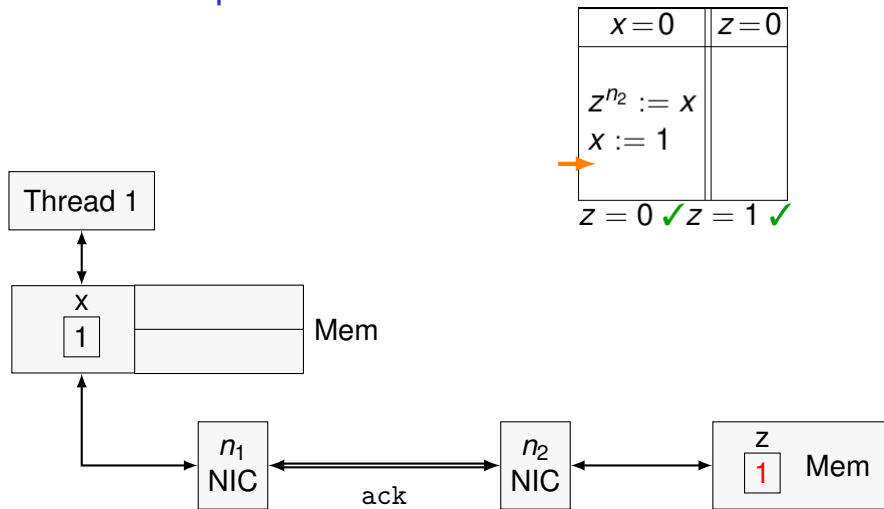
RDMA example 1/3



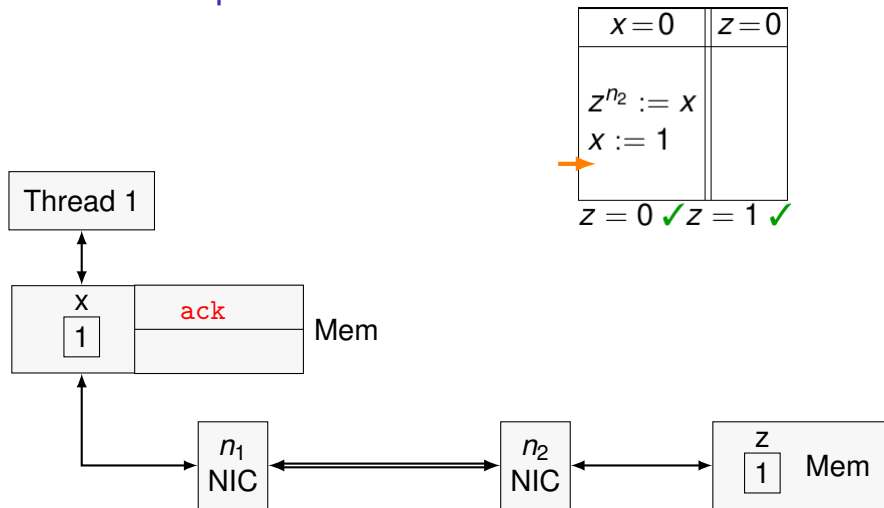
RDMA example 1/3



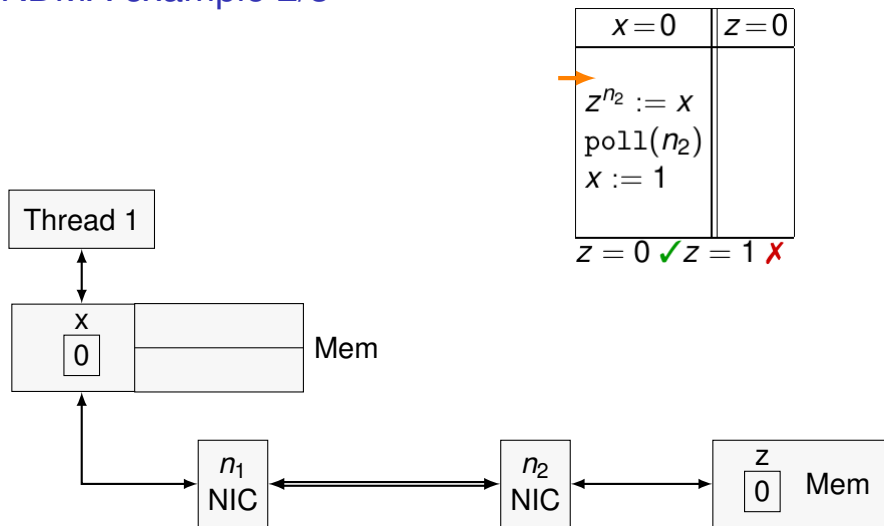
RDMA example 1/3



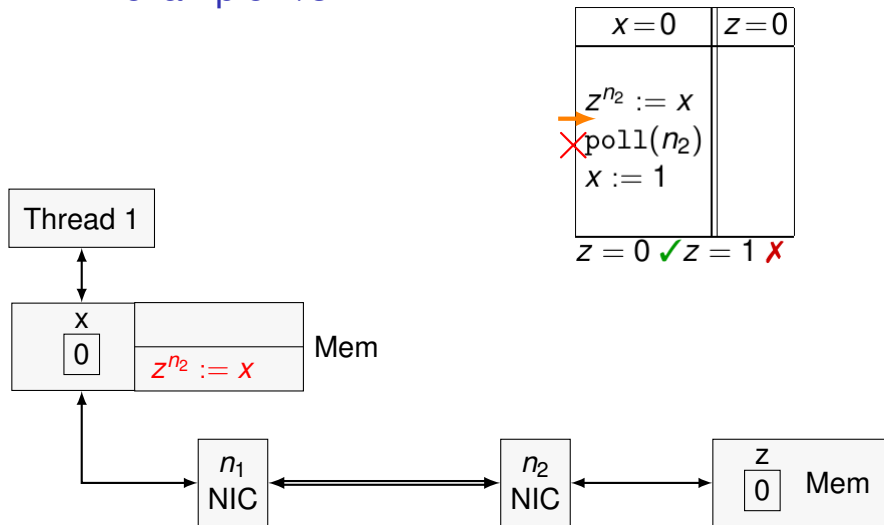
RDMA example 1/3



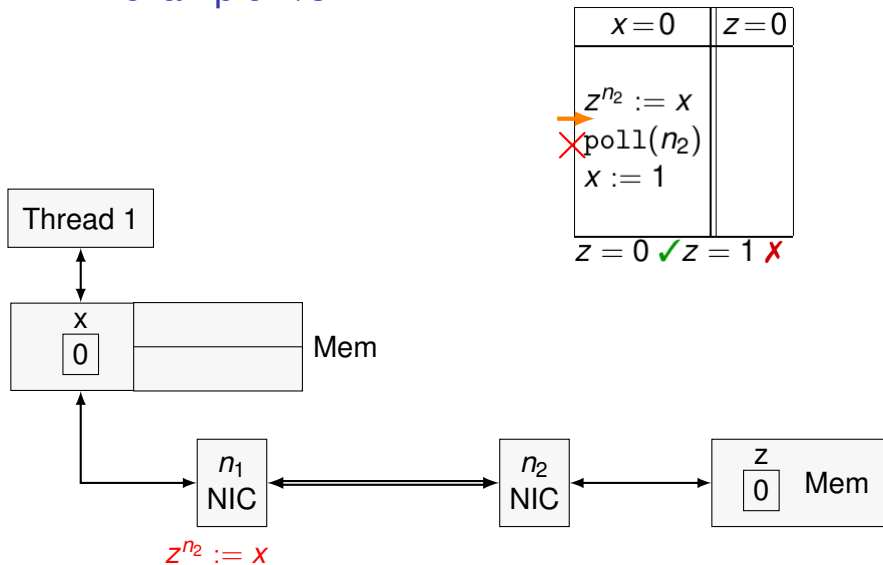
RDMA example 2/3



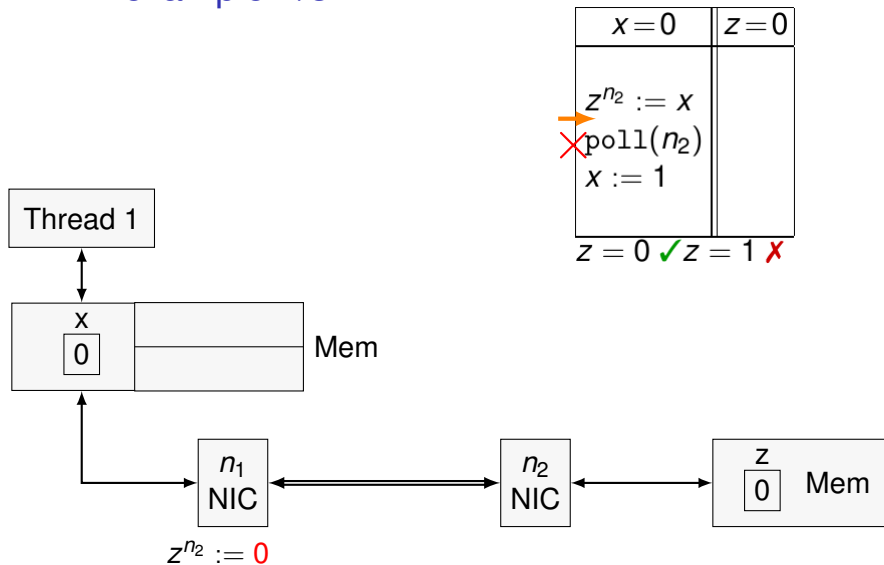
RDMA example 2/3



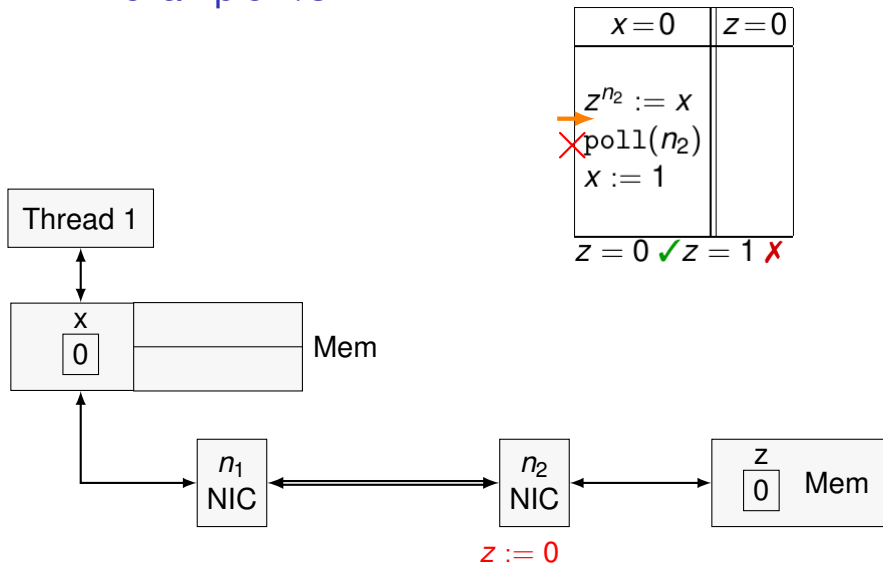
RDMA example 2/3



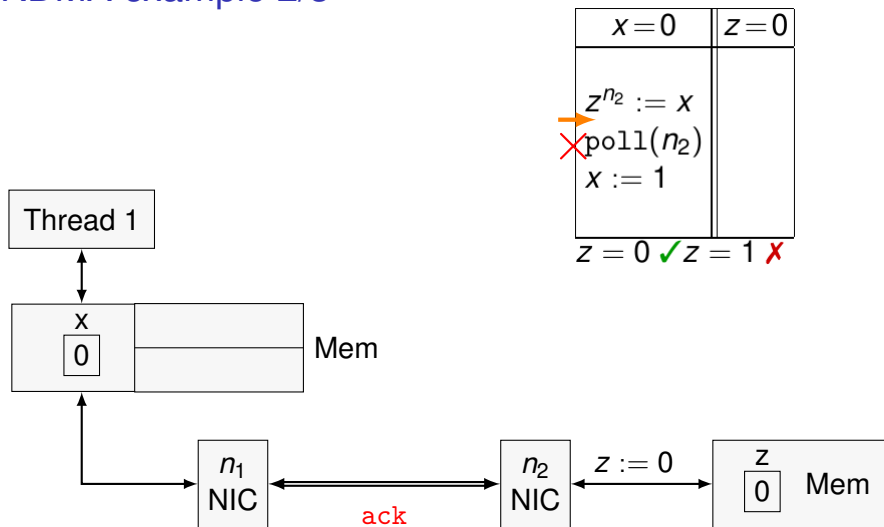
RDMA example 2/3



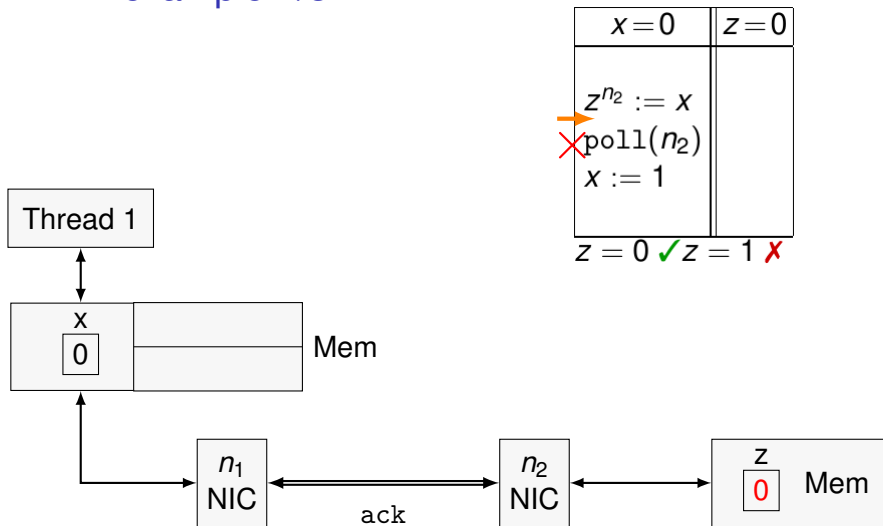
RDMA example 2/3



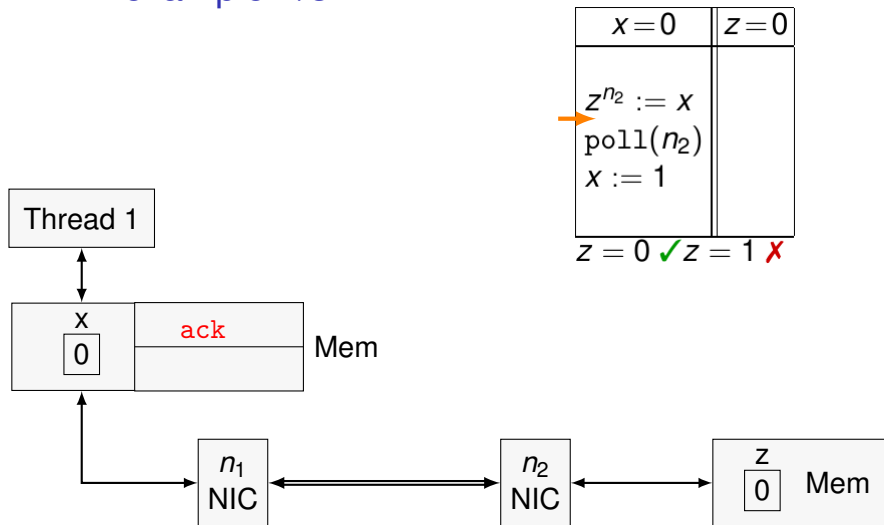
RDMA example 2/3



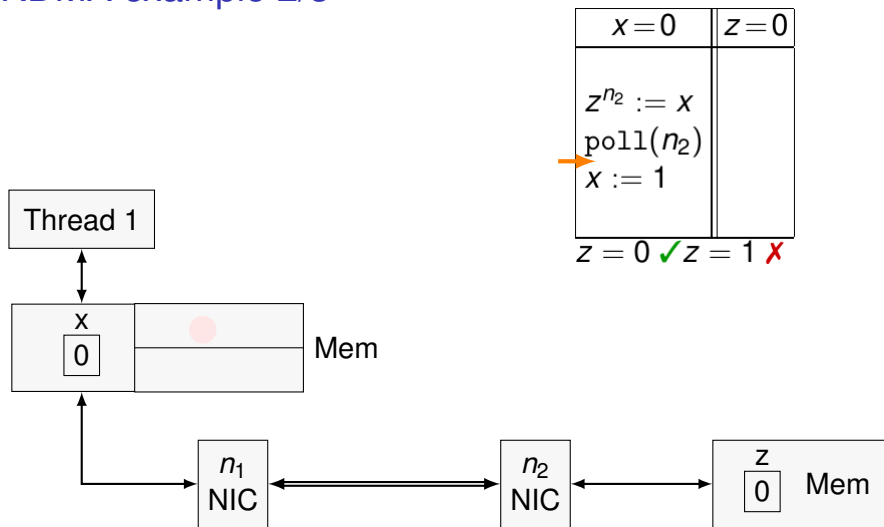
RDMA example 2/3



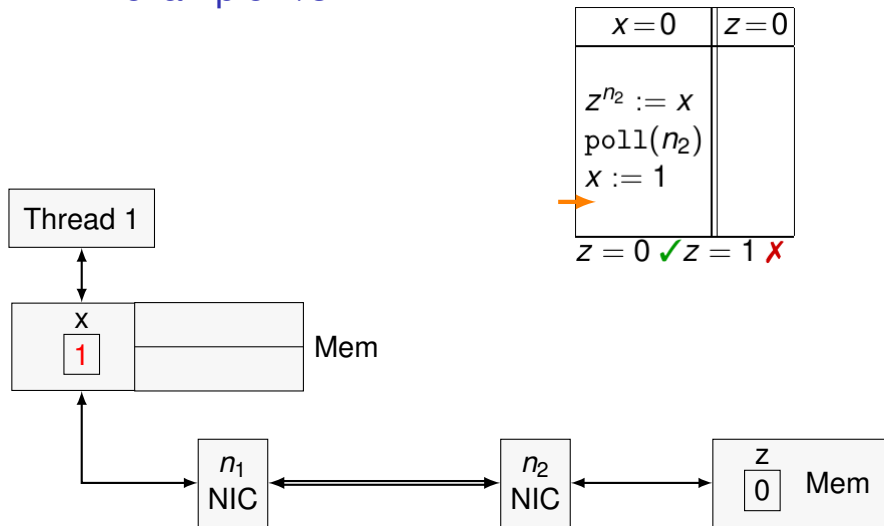
RDMA example 2/3



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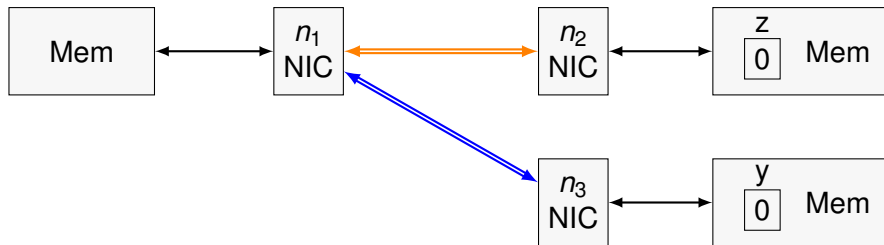
RDMA example 2/3



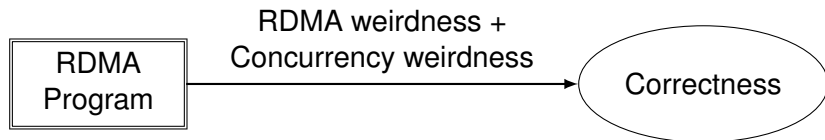
RDMA example 3/3

Independent paths are independent

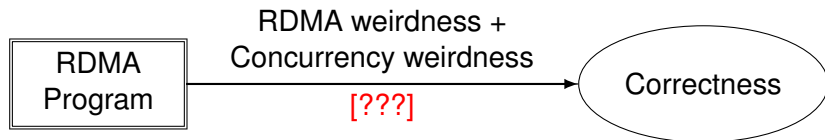
	$y = 0$	$z = 0$
$z^{n_3} := 1$ $y^{n_2} := 1$		



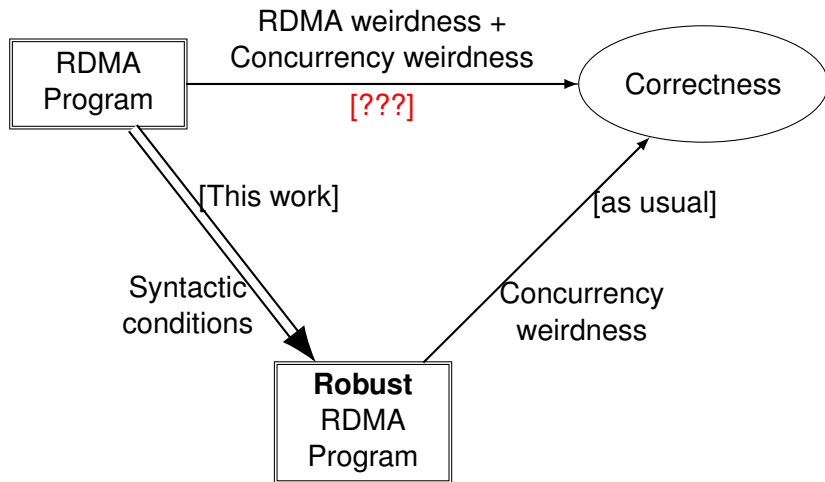
Formal Verification



Formal Verification



Formal Verification



SC and Robustness

Sequential Consistency (noun): Theoretical memory model where computers execute programs in order, line by line.

a.k.a.: **no fun allowed.**

Robust (adj.): Program whose outputs are achievable under SC.

a.k.a.: **it's not cheating unless you get caught!**

Preventing Reorderings

Result 1: existing RDMA programs can be made robust

Strategy: when in doubt, prevent reorderings

Preventing Reorderings

Result 1: existing RDMA programs can be made robust

	$y = 0$	$z = 0$
$z^{n_3} := 1$ $y^{n_2} := 1$		

ROBUST

Preventing Reorderings

Result 1: existing RDMA programs can be made robust

	$y = 0$	$z = 0$
$z^{n_3} := 1$ $y^{n_2} := 1$	$a := y$ $b := z^{n_3}$	

NOT ROBUST
 $(a, b) = (1, 0)$

Preventing Reorderings

Result 1: existing RDMA programs can be made robust

	$y = 0$	$z = 0$
$z^{n_3} := 1$ $\text{poll}(n_3)$ $y^{n_2} := 1$	$a := y$ $b := z^{n_3}$	

NOT ROBUST

$(a, b) = (1, 0)$

Preventing Reorderings

Result 1: existing RDMA programs can be made robust

$x = 0$	$y = 0$	$w, z = 0$
$z^{n_3} := 1$ $x := w^{n_3}$ $\text{poll}(n_3)$ $\text{poll}(n_3)$ $y^{n_2} := 1$	$a := y$ $b := z^{n_3}$	

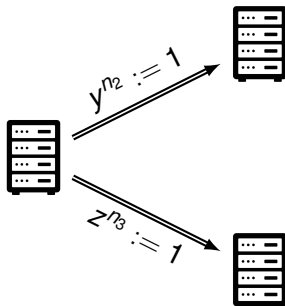
ROBUST

Tree topology

Result 2: programming guidelines for robustness
(a.k.a.: how to not get caught)

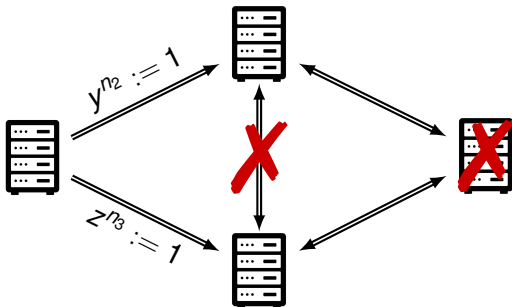
Tree topology

Result 2: programming guidelines for robustness



Tree topology

Result 2: programming guidelines for robustness



Condition 1: Tree topology

Maintaining plausible deniability

Result 2: programming guidelines for robustness

Standard RDMA usage:

$$x := y^n$$

$$\vdots$$
$$\vdots$$

$$f(x)$$

Maintaining plausible deniability

Result 2: programming guidelines for robustness

Standard RDMA usage:

$$x := y^n$$

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$$f(x)$$

NOT ROBUST

Maintaining plausible deniability

Result 2: programming guidelines for robustness

Standard RDMA usage:

$x := y^n$

$\vdots$

$\vdots$

`poll( $n$ )`

$f(x)$

ROBUST

Condition 2: Poll before reusing

Maintaining plausible deniability

Result 2: programming guidelines for robustness

Standard RDMA usage:

$$\left. \begin{array}{l} x := y^n \\ \vdots \\ z := 1 \\ \vdots \\ \text{poll}(n) \\ f(x) \end{array} \right\} \begin{array}{l} x \text{ can change} \\ \text{anytime} \end{array}$$

ROBUST

Maintaining plausible deniability

Result 2: programming guidelines for robustness

Standard RDMA usage:

$a := z$	$x := y^n$	} x can change anytime
	$\vdots$	
$b := x$	$z := 1$	
	$\vdots$	
	$\text{poll}(n)$	
	$f(x)$	

NOT ROBUST

$(a, b) = (1, 0)$

Maintaining plausible deniability

Result 2: programming guidelines for robustness

Standard RDMA usage:

	$x := y^n$	} x can change anytime
	$\vdots$	
$a := z$	$z := 1$	
$b := x$	$\vdots$	
	$\text{poll}(n)$	
	$f(x)$	

NOT ROBUST

$(a, b) = (1, 0)$

Condition 3: No snitches

(locations for RDMA operations are private)

Maintaining plausible deniability

Result 2: programming guidelines for robustness

Standard RDMA usage:

	$x := y^n$	}	x can change anytime
	$\vdots$		
	$z := 1$		
$a := z$	$\vdots$		
$b := x'$	$\text{poll}(n)$		
	$x' := x$		
	$f(x)$		

ROBUST

Theorem

Result 2: programming guidelines for robustness

- Tree topology
- Poll before reusing
- No snitches

} $\Rightarrow$ Robustness

Conclusion

RDMA

- Low-latency data-transfer protocol for HPC/databases
- Recently becoming widespread
- Very weak and unintuitive semantics

Our work: RDMA robustness (assuming x86-TSO for the CPU)

- Existing RDMA programs can be made robust (and probably slow)
- With tree topology, can write robust RDMA programs for cheap
- Else... good luck for your formal verification